

Graph Properties of Graph Associahedra

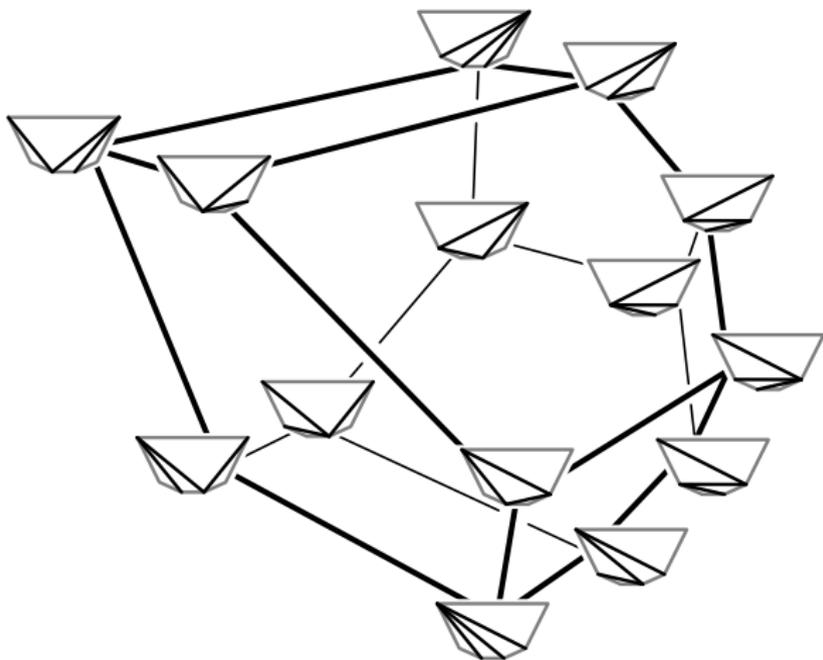
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joint work with **Vincent Pilaud** (CNRS, LIX Polytechnique)

March 24th, 2014

Definition

An *associahedron* is a polytope whose graph is the flip graph of triangulations of a convex polygon.



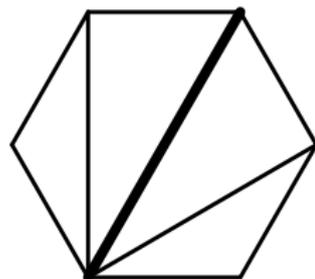
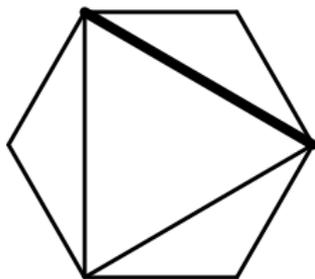
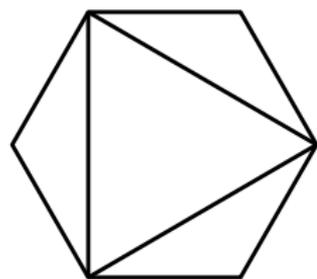
Faces \leftrightarrow dissections of the polygon

Focus on graphs

Flip graph on the triangulations of the polygon:

Vertices: *triangulations*

Edges: *flips*

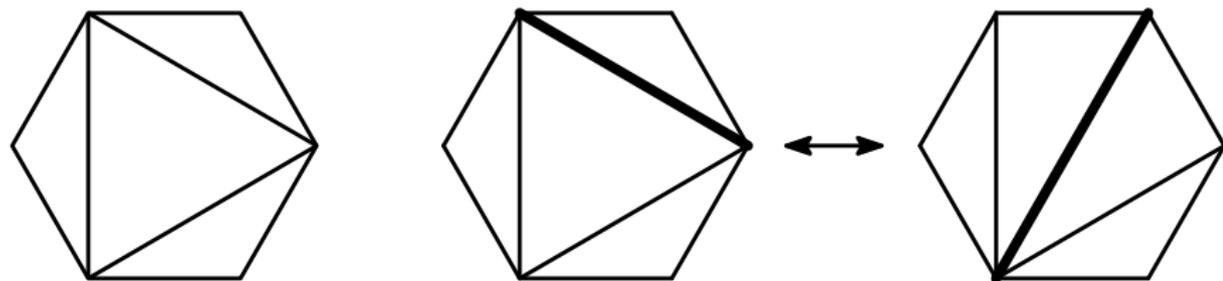


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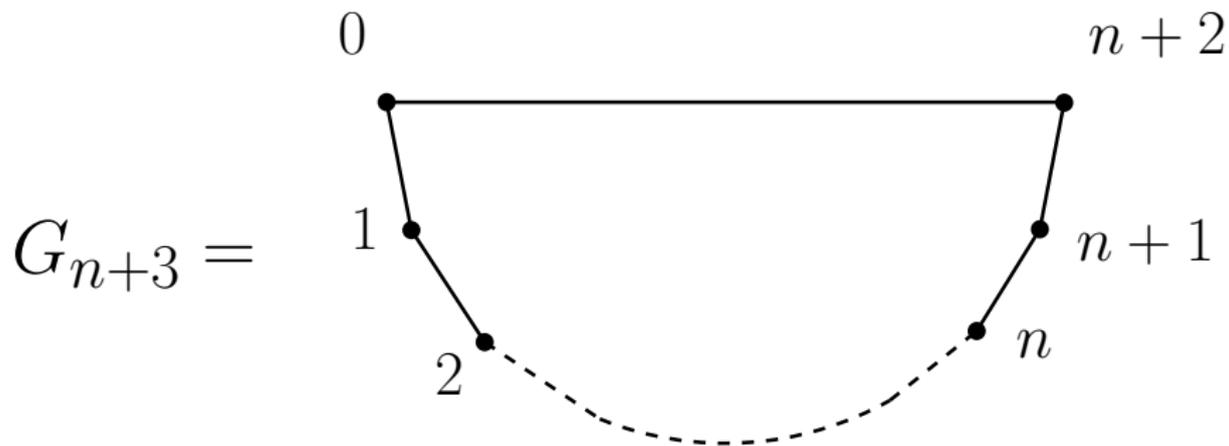
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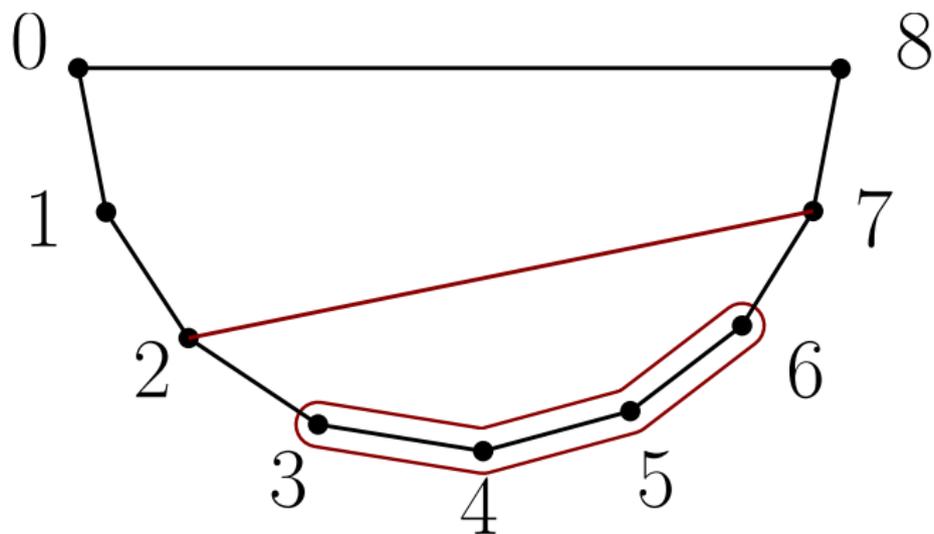
n diagonals \Rightarrow the flip graph is n -regular.

Useful configuration (Loday's)



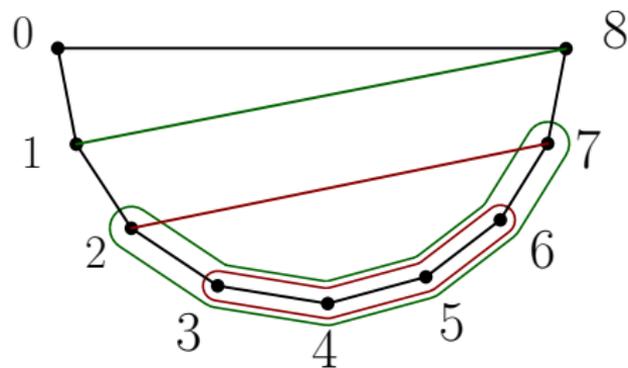
Graph point of view

$\{\text{diagonals of } G_{n+3}\} \longleftrightarrow \{\text{strict subpaths of the path } [n+1]\}$

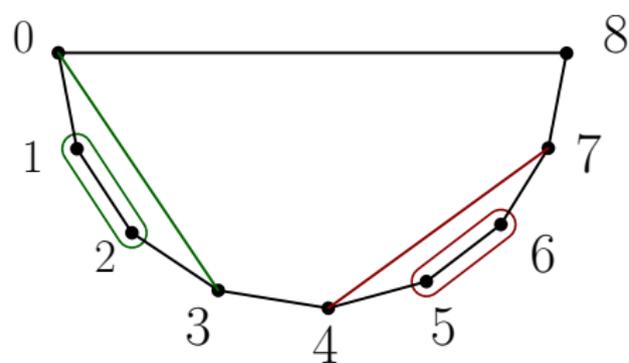


Non-crossing diagonals

Two ways to be non-crossing in Loday's configuration:



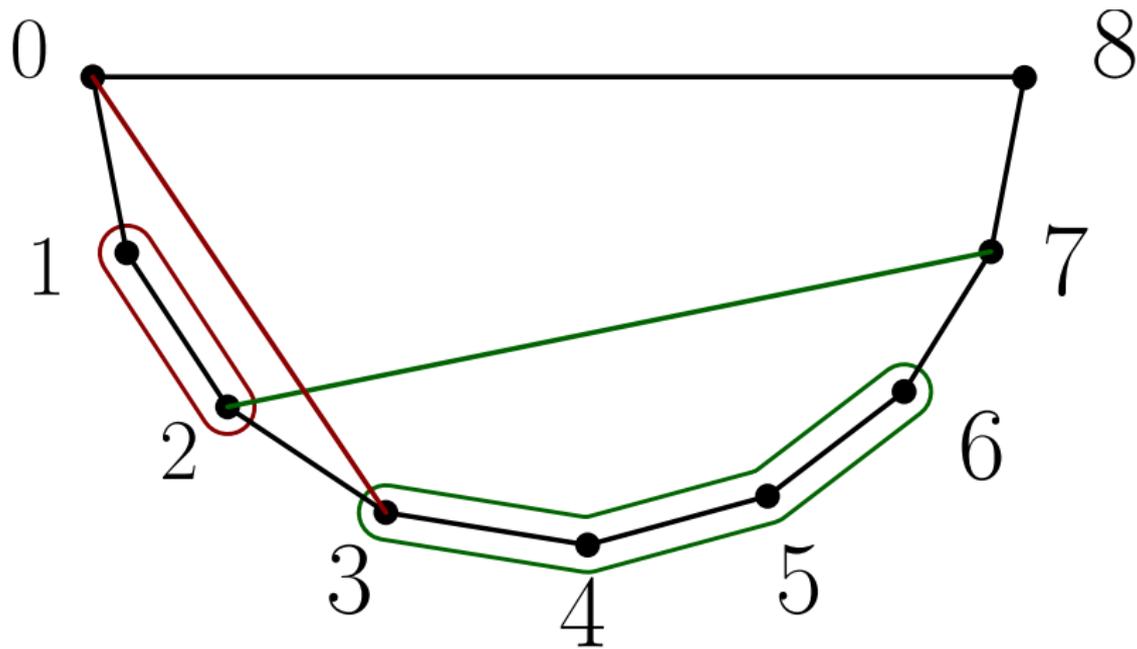
nested subpaths



non-adjacent subpaths

Pay attention to the second case:

The right condition is indeed *non-adjacent*, disjoint is not enough!



Now do it on graphs

$G = (V, E)$ a (connected) graph.

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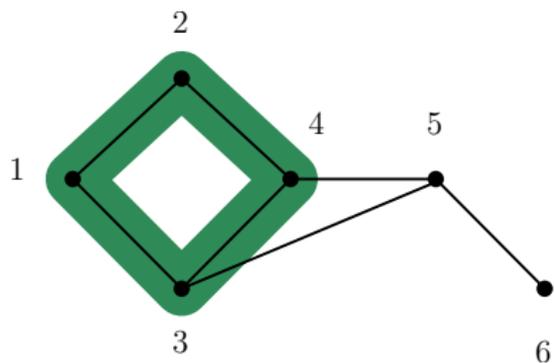
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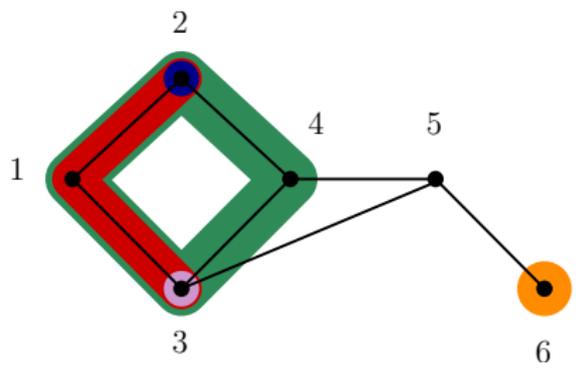
Definition

- A **tube** of G is a proper subset $t \subseteq V$ inducing a connected subgraph of G ;
- t and t' are **compatible** if they are nested or non-adjacent;
- A **tubing** of G is a set of pairwise compatible tubes of G .



A tube

(generalizes a diagonal)



A maximal tubing

(generalizes a triangulation)

Graph associahedra

The simplicial complex of tubings is spherical

Graph associahedra

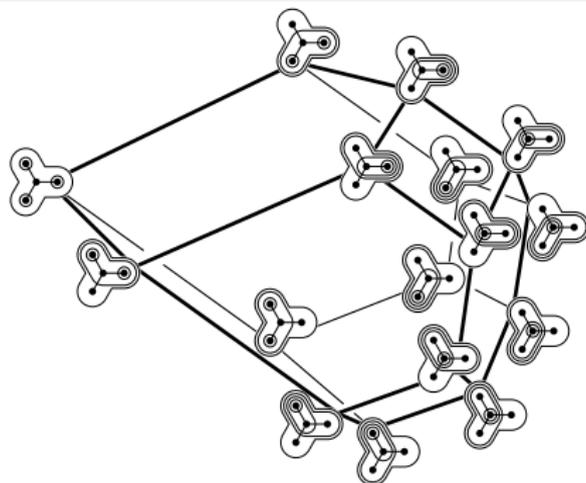
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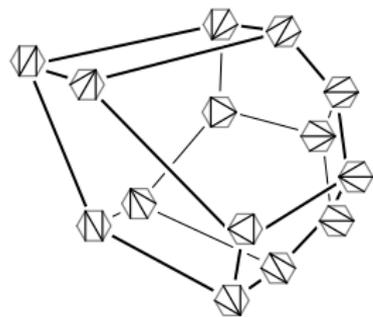
Theorem (Carr-Devadoss '06)

*There exists a polytope called **graph associahedron** of G , denoted \mathbf{Asso}_G , whose graph is this flip graph.*

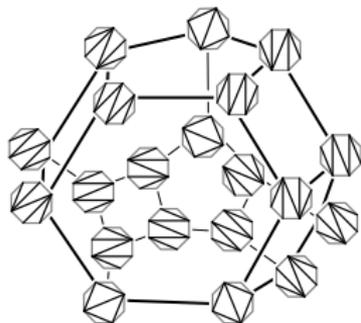


Faces \leftrightarrow tubings of G .

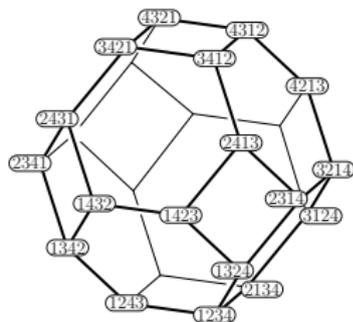
Classical polytopes...



The associahedron



The cyclohedron



The permutahedron

Hamiltonicity of flip graphs

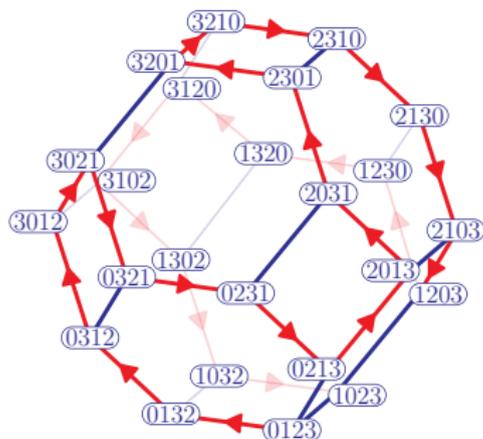
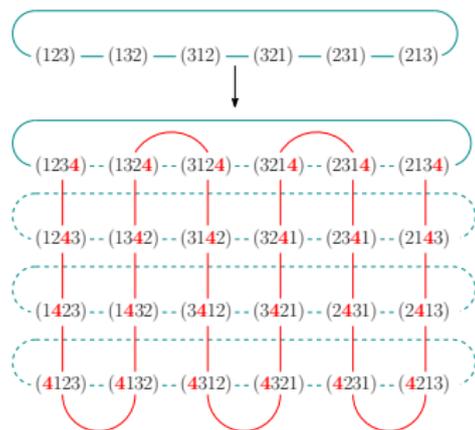
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The n -dimensional permutahedron is hamiltonian for $n \geq 2$.

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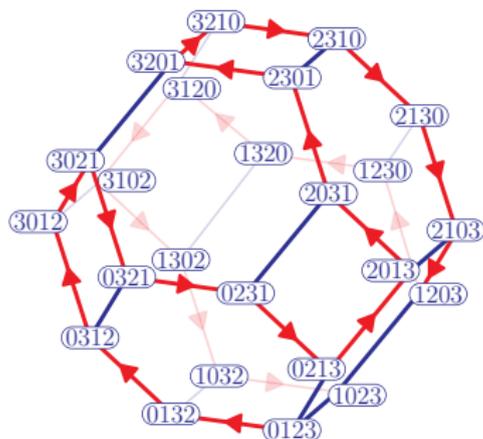
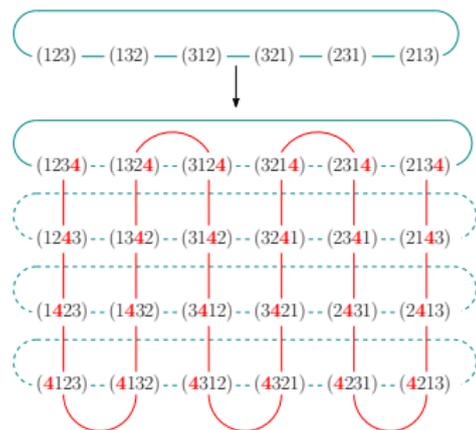
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Hamiltonicity of flip graphs

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Theorem (Lucas 87, Hurtado-Noy '99)

The n -dimensional associahedron is hamiltonian for $n \geq 2$.

Hamiltonicity

Theorem (M.-Pilaud '14⁺)

Any graph associahedron is hamiltonian.

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- Truncation hyperplanes correspond to tubes.

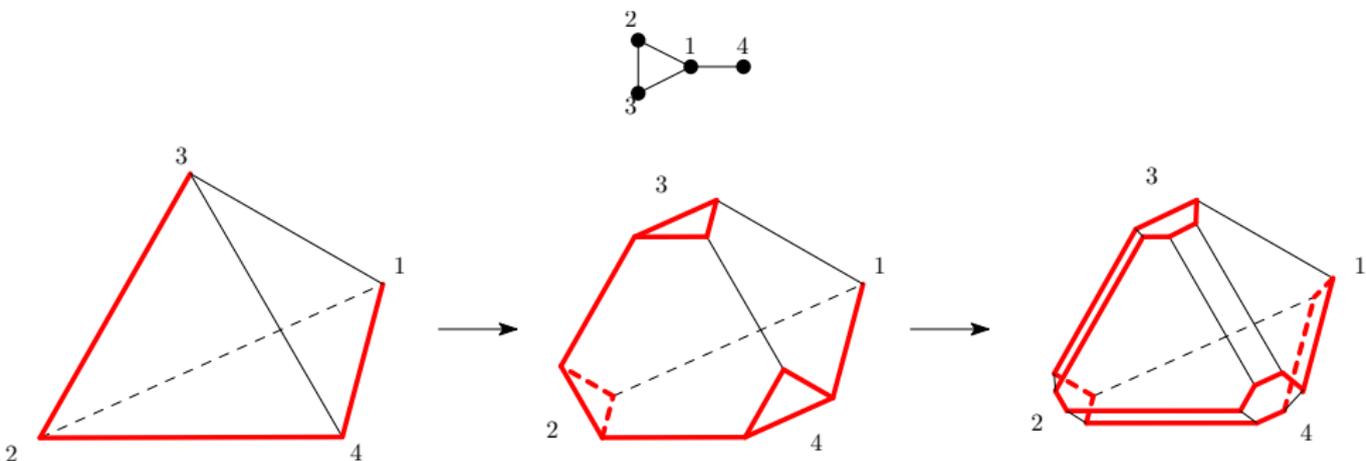
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Diameter of flip graphs

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Theorem (Sleator-Trajan-Thurston '88, Pournin '12)

The diameter of the n -dimensional associahedron is $2n - 4$ for $n \geq 10$.

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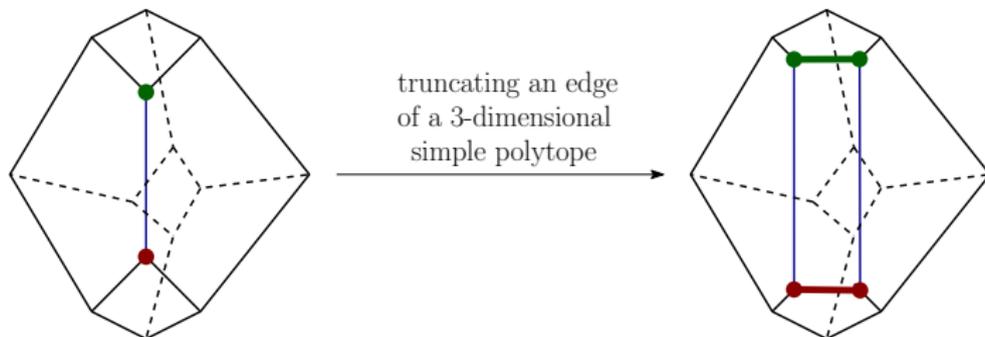
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→ Truncating \iff replacing vertices by complete graphs.



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- Pournin's result for the classical associahedron.

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THANK YOU FOR
YOUR ENTHUSIASTIC
ATTENTION !